

# Planes, Hyperplanes, and Beyond

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## Summary

Suppose that you have a cake and are allowed to make ten straight-line slices. What is the greatest number of pieces you can produce? What if the slices have to be symmetric — or if the cake is four-dimensional? How can we possibly see what it looks like to slice space into pieces using lines, planes, or hyperplanes? Many of these questions have beautiful answers that can be revealed using unexpected, yet essentially simple mathematical techniques. Better yet, the seemingly abstract study of hyperplane arrangements has many surprising practical applications, ranging from optimization problems, to the theory of networks, to how a group of cars can find parking spots.

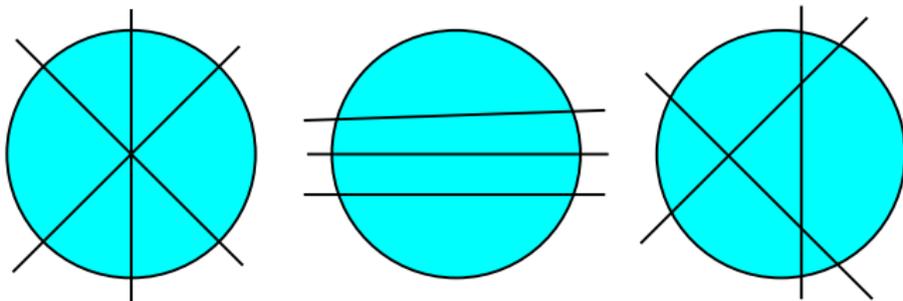
# Organization

- Part 1: Cutting a cake into as many pieces as possible
- Part 2: Symmetric cake-cutting
- Part 3: Parking cars, planting trees, scoring with handicaps, and what all that has to do with cake-cutting

# Part 1: How Many Pieces of Cake?

# The Cake-Cutting Problem

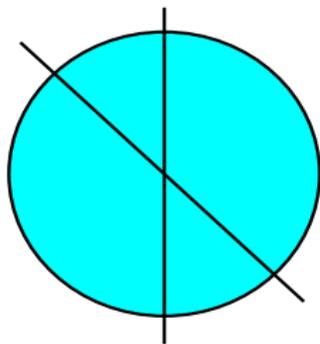
What is the greatest number of pieces that a cake can be cut into with a given number of cuts?



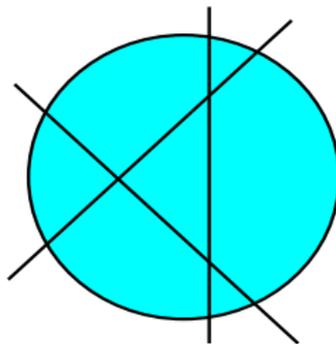
- ▶ The cuts must be **straight lines** and must go **all the way through** the cake.
- ▶ The sizes and shapes of the pieces **don't matter**.
- ▶ For the moment, we'll focus on **2-dimensional** cakes (think of them as pancakes).

## Solutions with 2, 3 or 4 Cuts

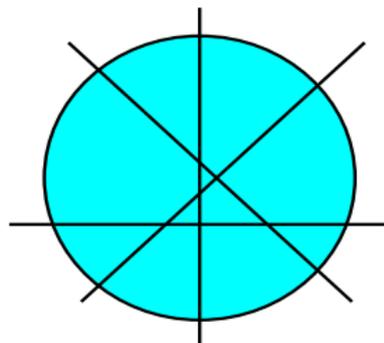
Let's write  $P_2(N)$  for the maximum number of pieces obtainable using  $N$  cuts. (The 2 stands for dimension.)



2 cuts:  
 $P_2(2) = 4$



3 cuts:  
 $P_2(3) = 7$



4 cuts:  
 $P_2(4) = 11$

## Solutions with $N$ Cuts

Cuts $N$	Pieces $P_2(N)$
1	2
2	4
3	7
4	11
5	16
6	22

Cuts $N$	Pieces $P_2(N)$
7	29
8	37
9	46
10	56
...	...
100	5051

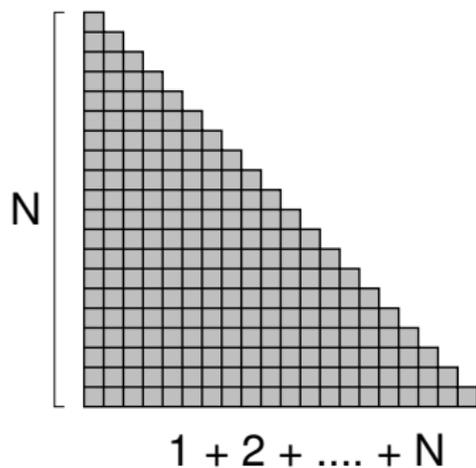
Do you see the pattern?

# The Pattern

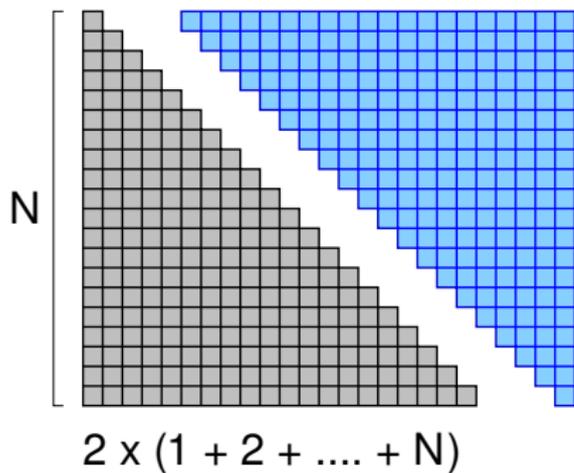
$N$	$P_2(N)$			
0	1			
1	2	=	1 + 1	
2	4	=	2 + 2	= 1 + 1 + 2
3	7	=	4 + 3	= 1 + 1 + 2 + 3
4	11	=	7 + 4	= 1 + 1 + 2 + 3 + 4
5	16	=	11 + 5	= 1 + 1 + 2 + 3 + 4 + 5
...	...	=	...	= ...

- ▶ How do we **prove** that the pattern works for **every**  $N$ ?
- ▶ What does  $1 + 2 + \dots + N$  equal anyway?

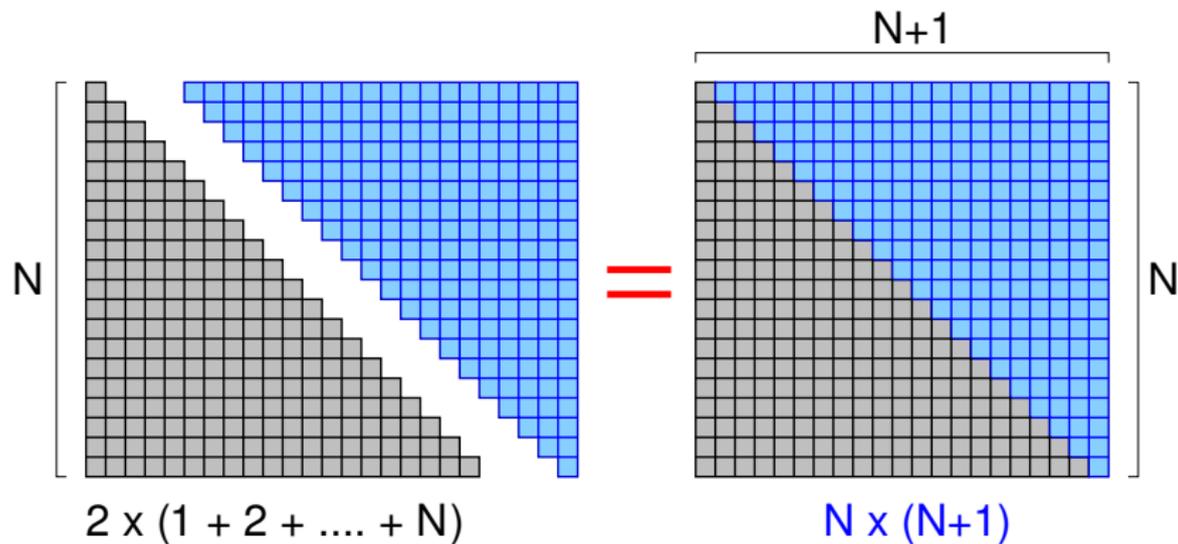
# The Staircase Theorem



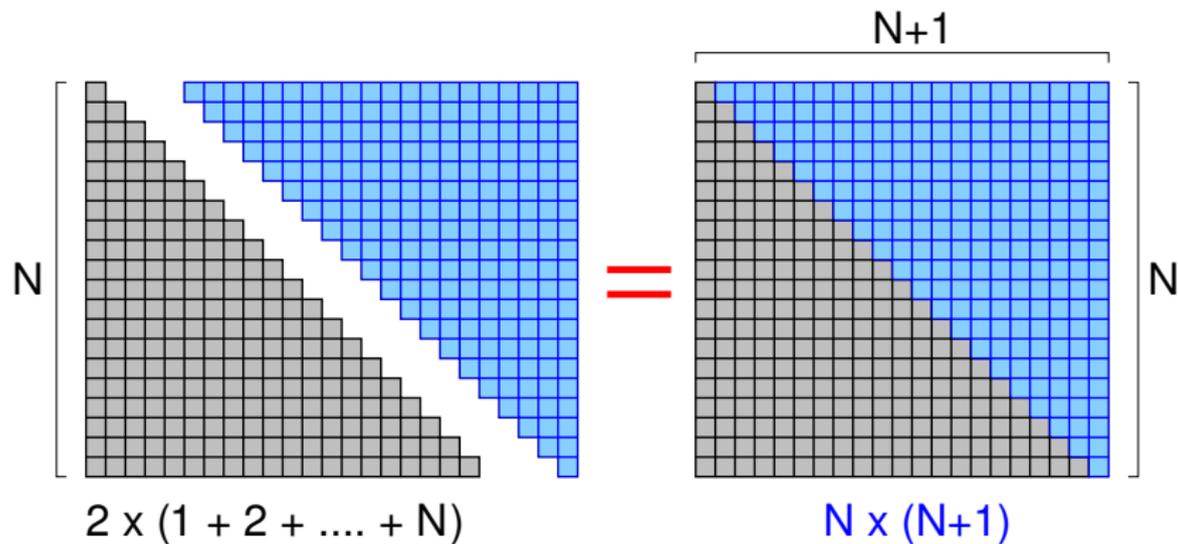
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# The Staircase Theorem



$$1 + 2 + \dots + N = N(N+1) / 2$$

# The Pattern

$N$	$P_2(N)$			
0	1			
1	2	=	$1 + 1$	
2	4	=	$2 + 2$	= $1 + 1 + 2$
3	7	=	$4 + 3$	= $1 + 1 + 2 + 3$
4	11	=	$7 + 4$	= $1 + 1 + 2 + 3 + 4$
5	16	=	$11 + 5$	= $1 + 1 + 2 + 3 + 4 + 5$
...	...	...	...	...

By the Staircase Theorem, we can **conjecture** that

$$P_2(N) = 1 + (1 + 2 + \cdots + N) = 1 + \frac{N(N+1)}{2}.$$

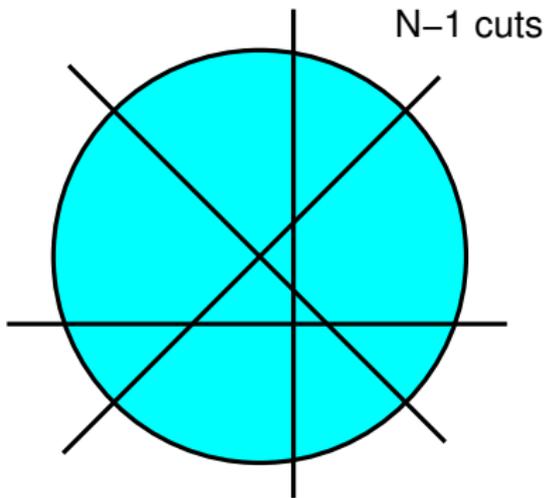
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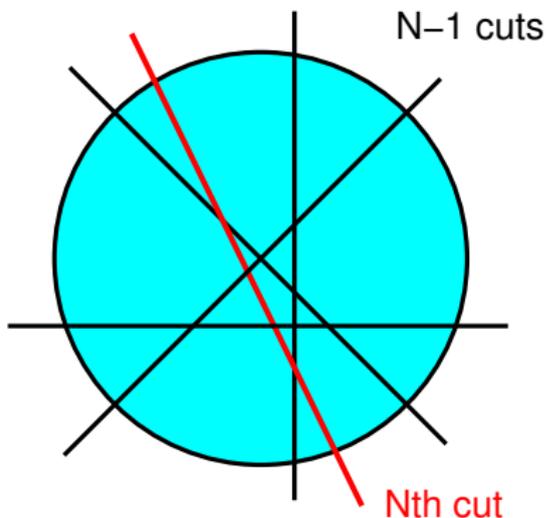
- ▶ First cut the pancake into  $P_2(N - 1)$  pieces using  $N - 1$  cuts.



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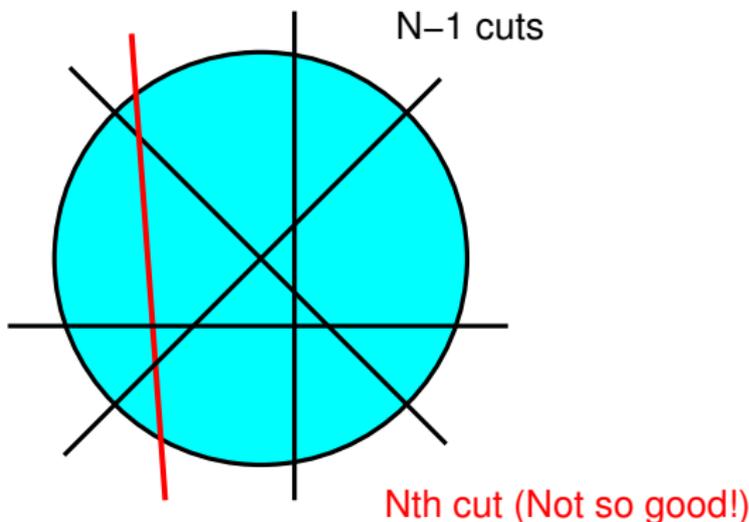
- ▶ First cut the pancake into  $P_2(N - 1)$  pieces using  $N - 1$  cuts.
- ▶ Now make the  $N$ th cut, hitting as many pieces as possible.



# Maximizing the Number of Pieces

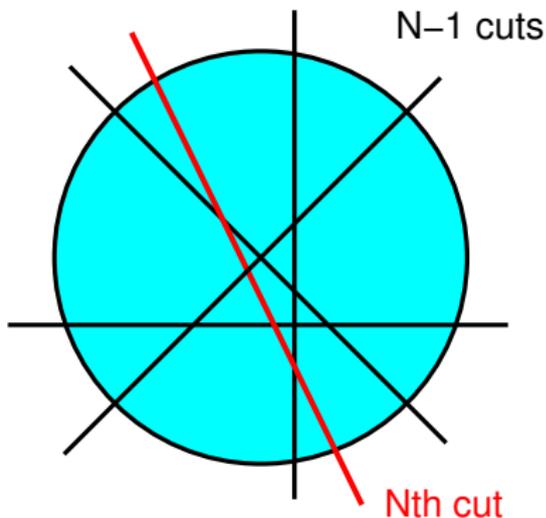
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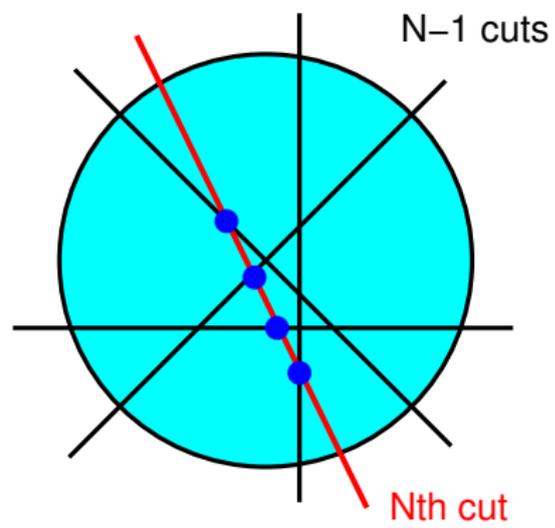
# Maximizing the Number of Pieces

**Key observation:** The **number of pieces** subdivided by the  $N$ th cut equals **one more than the number of previous cuts it meets**.



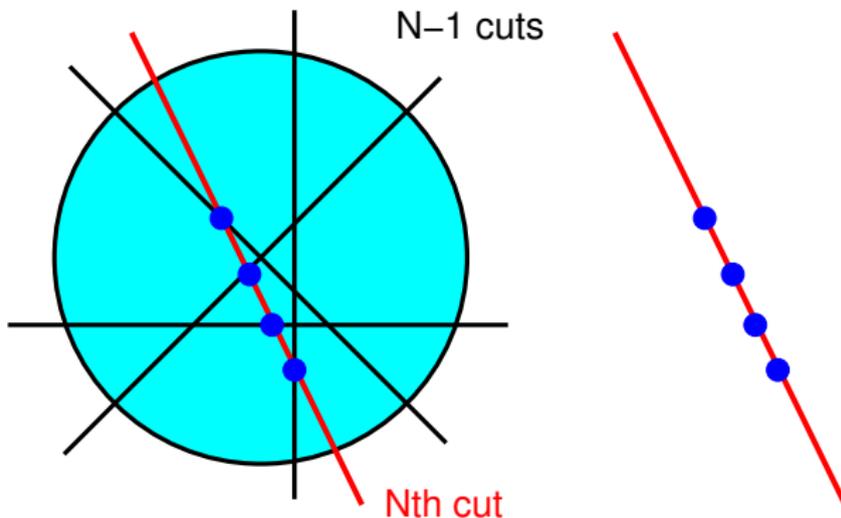
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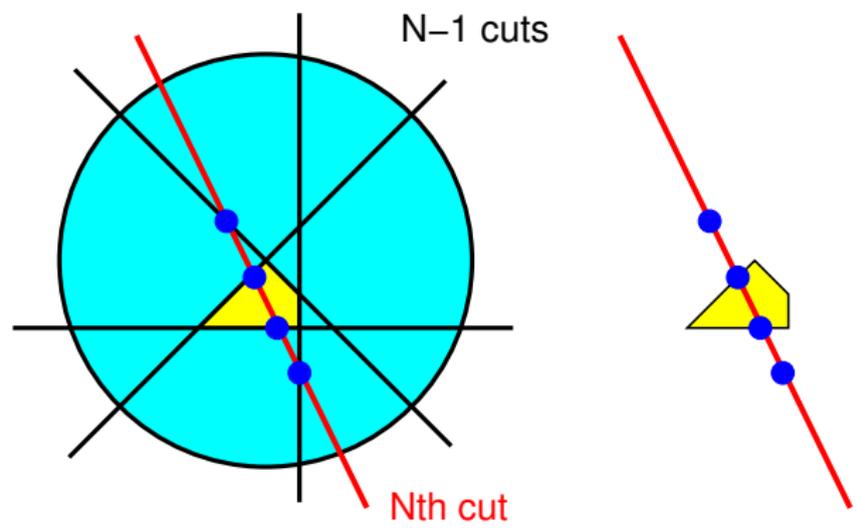
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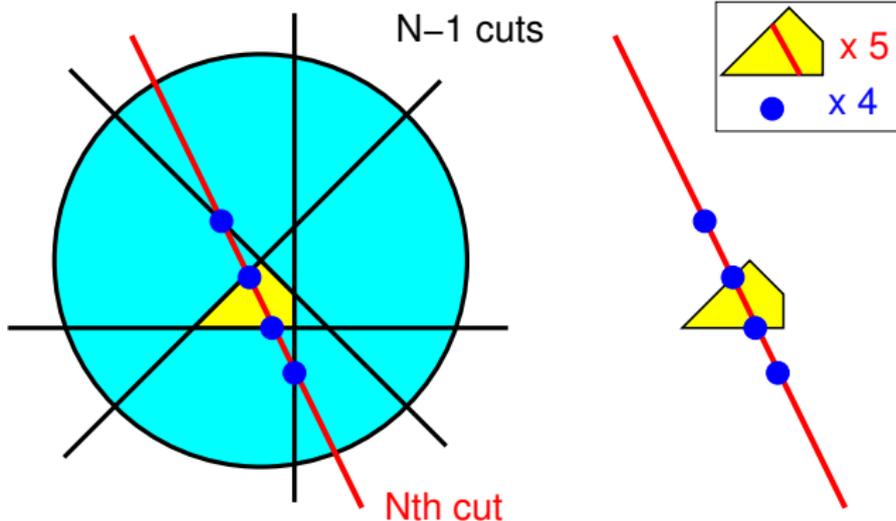
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# Maximizing the Number of Pieces

If we make sure that

- ▶ every pair of cuts meets in some point, and
- ▶ no more than two cuts meet at any point,

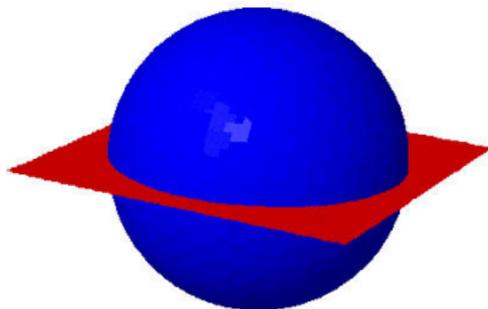
then the  $N^{\text{th}}$  cut will meet each of the previous  $N - 1$  cuts, and therefore will make  $N$  new pieces.

Since the original pancake had one piece, we have **proved** that

$$P_2(N) = 1 + (1 + 2 + \cdots + N) = 1 + \frac{N(N + 1)}{2}.$$

## From 2D to 3D

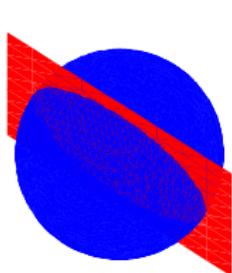
What about 3-dimensional cakes?



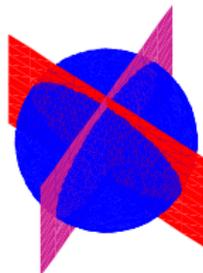
A cut in 3-dimensional space means a **plane**, not a line.

## From 2D to 3D

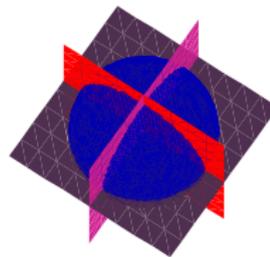
Let's write  $P_3(N)$  for the maximum number of pieces obtainable from a 3-dimensional cake with  $N$  cuts.



$$P_3(1) = 2$$



$$P_3(2) = 4$$

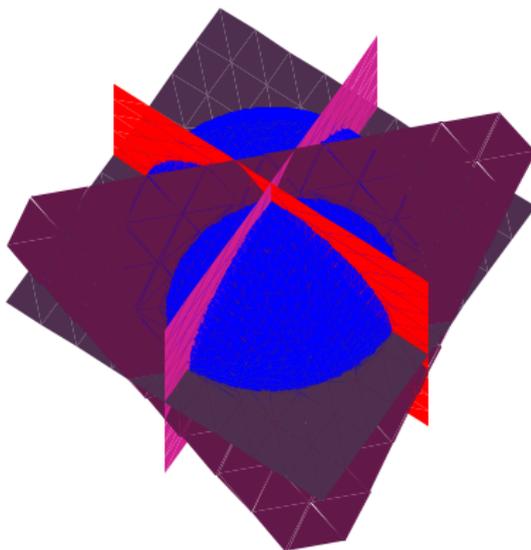


$$P_3(3) = 8$$

Compare 2D:  $P(1) = 2$ ,  $P(2) = 4$ ,  $P(3) = 7$ .

$$P_3(4) = 15$$

With four planes, we can make 15 pieces (though only 14 are visible from the outside).



# From 2D to 3D

$N$	0	1	2	3	4	5	6	7	8
$P_2(N)$	1	2	4	7	11	16	22	29	37
$P_3(N)$	1	2	4	8	15	26	42	64	93

Do you see the pattern?

## From 2D to 3D

$N$	0	1	2	3	4	5	6	7	8
$P_2(N)$	1	2	4	7	<b>11</b>	16	22	29	37
$P_3(N)$	1	2	4	8	<b>15</b>	<b>26</b>	42	64	93

The pattern is

$$P_3(N) = P_3(N - 1) + P_2(N - 1).$$

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The pattern is

$$P_3(N) = P_3(N-1) + P_2(N-1).$$

(In fact  $P_3(N) = \frac{N^3+5N+6}{6}$  — but the pattern is more important than this formula!)

# Pancakes, Cakes and Hypercakes

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In general, if you have a  $d$ -dimensional cake and you can make  $N$  cuts, how many pieces can you make? (Call this number  $P_d(N)$ .)

- ▶ We already know the answers for  $d = 2$  and  $d = 3$ .
- ▶ For  $d = 1$ :  $N$  cuts give  $N + 1$  pieces.
- ▶ For any  $d$ : 0 cuts give 1 piece, 1 cut gives 2 pieces.

## Pancakes, Cakes and Beyond

- ▶ Each number is the sum of the numbers immediately “west” ( $\leftarrow$ ) and “northwest” ( $\nearrow$ ).
- ▶ Formula:  $P_d(N) = P_d(N-1) + P_{d-1}(N-1)$ .

	N								
	0	1	2	3	4	5	6	7	8
$P_1(N)$	1	2	3	4	5	6	7	8	9
$P_2(N)$	1	2	4	7	11	16	22	29	37
$P_3(N)$	1	2	4	8	15	26	42	64	93
$P_4(N)$	1	2	4	8	16	31	57	99	163
$P_5(N)$	1	2	4	8	16	32	63	120	219
...	...								

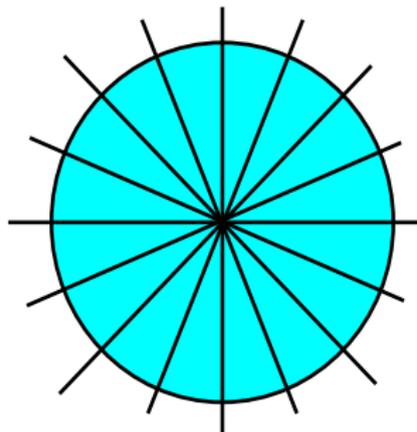
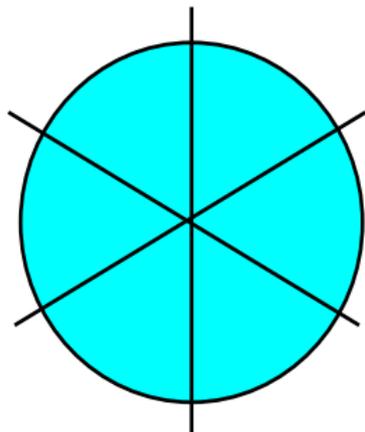
# Pancakes, Cakes and Beyond

**Theme:** Understanding patterns in dimensions we **can** see enables us to understand dimensions we **can't** see.

## Part 2: Symmetric Cake-Cutting

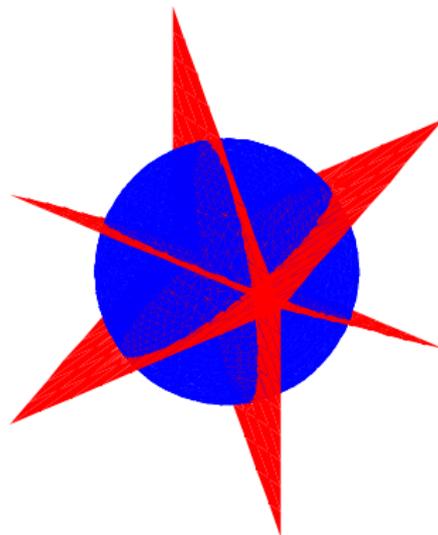
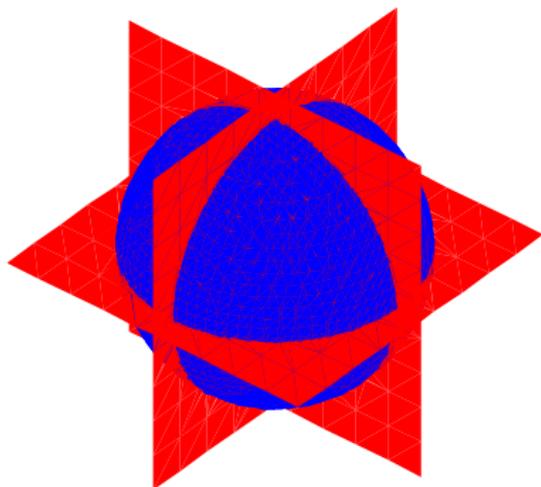
# Symmetric Cake-Cutting

What are the possible ways to cut a perfectly round cake so that **all pieces are congruent** (i.e., geometrically the same)?



# Symmetric Cake-Cutting

What are the possible ways to cut a perfectly round cake so that **all pieces are congruent** (i.e., geometrically the same)?



## Refresher: $N$ -Dimensional Algebra

Lines in 2-dimensional space have equations like

$$x = y, \quad x = 0, \quad x + 2y = 4.$$

Planes in 3-dimensional space have equations like

$$x = y, \quad x = z, \quad x = 0, \quad x + 3y + 2z = 1.$$

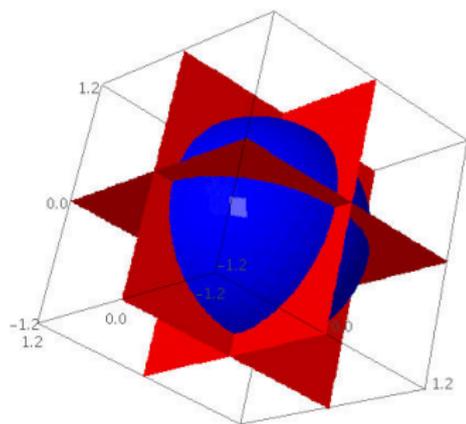
Hyperplanes in 4-dimensional space have equations like

$$x + y = z, \quad w = 0, \quad 3w - 2x + 7y + 2z = 2012.$$

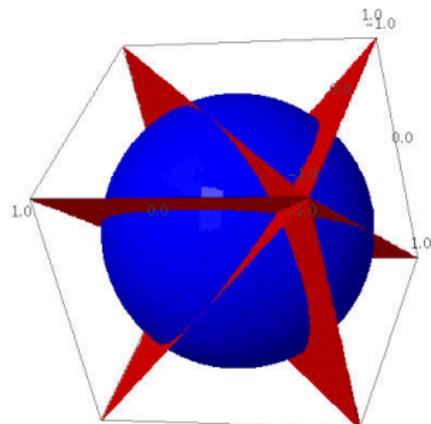
The two sides of a hyperplane are given by **inequalities**. For example, the plane  $x = z$  cuts 3D-space into the two pieces

$$x < z, \quad z < x.$$

# Symmetric Cake-Cutting



$$x = 0, y = 0, z = 0$$



$$x = y, x = z, y = z$$

# Symmetric Cake-Cutting in Higher Dimensions

**Question:** If we can cut up a **3-dimensional** sphere into congruent pieces using the **planes** defined by the equations

$$\boxed{x = 0, y = 0, z = 0} \quad \text{or} \quad \boxed{x = y, x = z, y = z}$$

then what happens if we cut up a **4-dimensional** sphere into pieces using the **hyperplanes**

$$\boxed{\begin{array}{ll} w = 0 & x = 0 \\ y = 0 & z = 0 \end{array}} \quad \text{or} \quad \boxed{\begin{array}{lll} w = x & w = y & w = z \\ x = y & x = z & y = z \end{array}} \quad ?$$

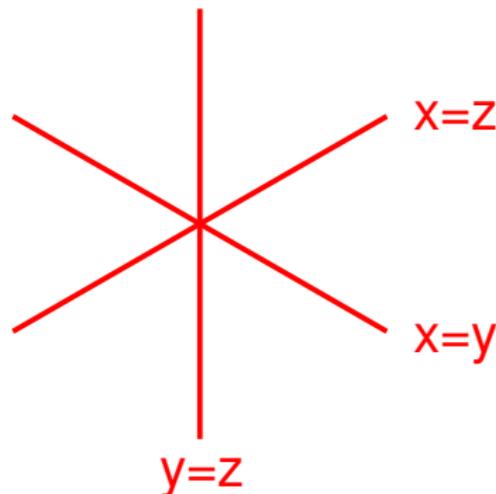
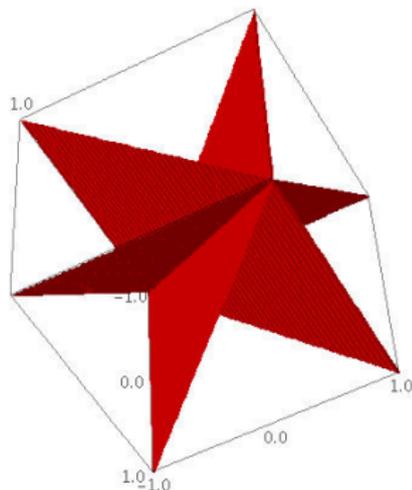
# Symmetric Cake-Cutting in Higher Dimensions

Some tools for visualizing 4-dimensional space:

- ▶ Work by **analogy**: understanding low-dimensional space can help us understand higher dimensions
- ▶ **Project** into lower dimension to make visualization easier
- ▶ **Reexpress** high-dimensional problems mathematically

# The Braid Arrangement

The arrangement of planes  $x = y$ ,  $x = z$ ,  $y = z$  is called the *3-dimensional braid arrangement* (**Braid3** for short).



Projecting from 3D to 2D makes the diagrams simpler, and retains both the number and symmetry of the regions.

## Regions Between The Planes of Braid<sub>3</sub>

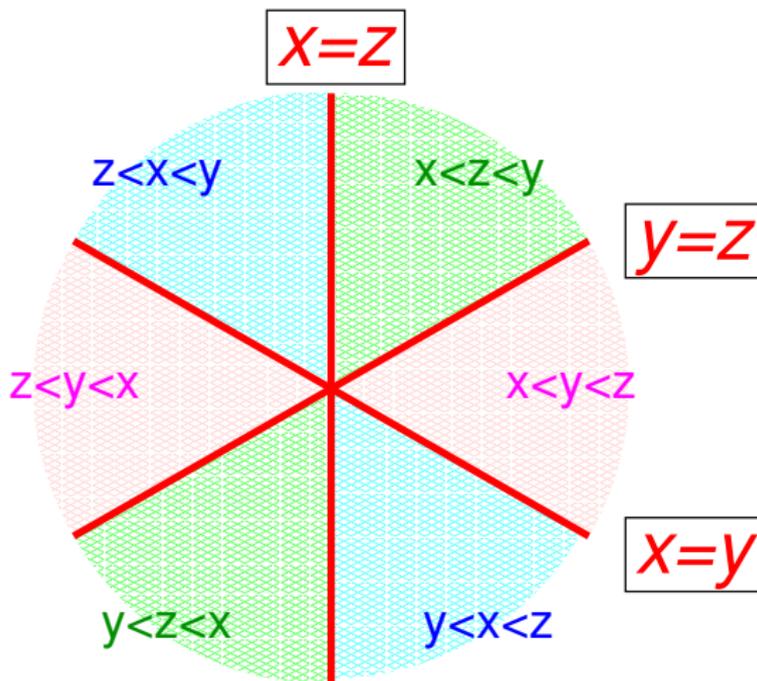
Each region of **Braid<sub>3</sub>** is on one side of each of the planes  $x = y$ ,  $x = z$ ,  $y = z$ . Therefore,

- ▶ either  $x < y$  or  $y < x$ ,
- ▶ either  $x < z$  or  $z < x$ , and
- ▶ either  $y < z$  or  $z < y$ .

Each region can be completely specified by the **order** of the three coordinates  $x, y, z$ . There are six possibilities:

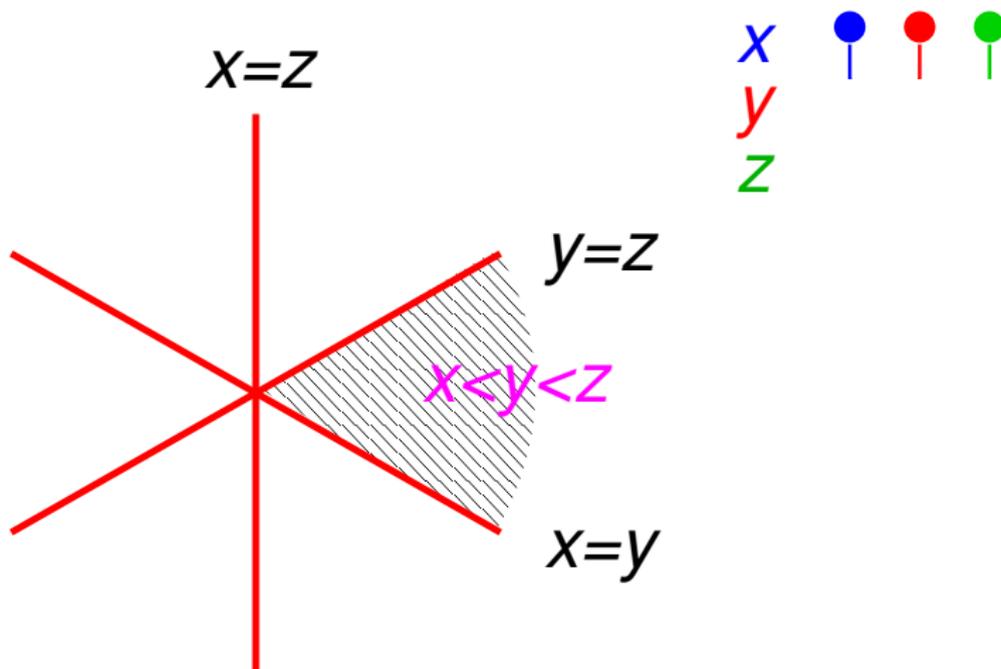
$$\begin{array}{lll} x < y < z & y < x < z & z < x < y \\ x < z < y & y < z < x & z < y < x \end{array}$$

# Regions of Braid<sub>3</sub>



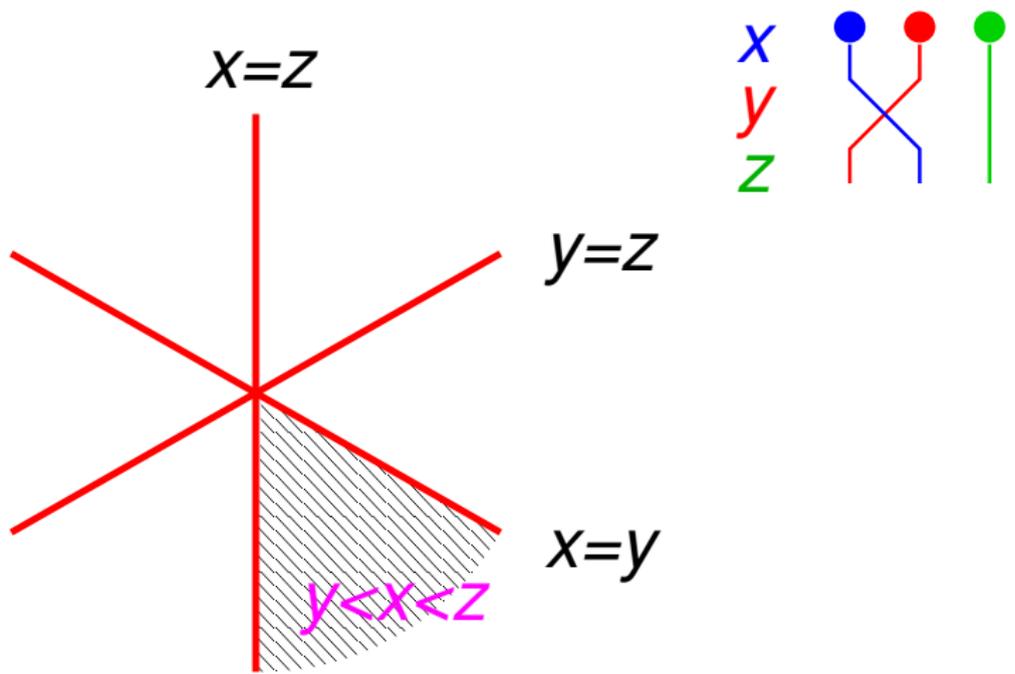
## Why “Braid”?

Crossing a border corresponds to reversing one inequality.



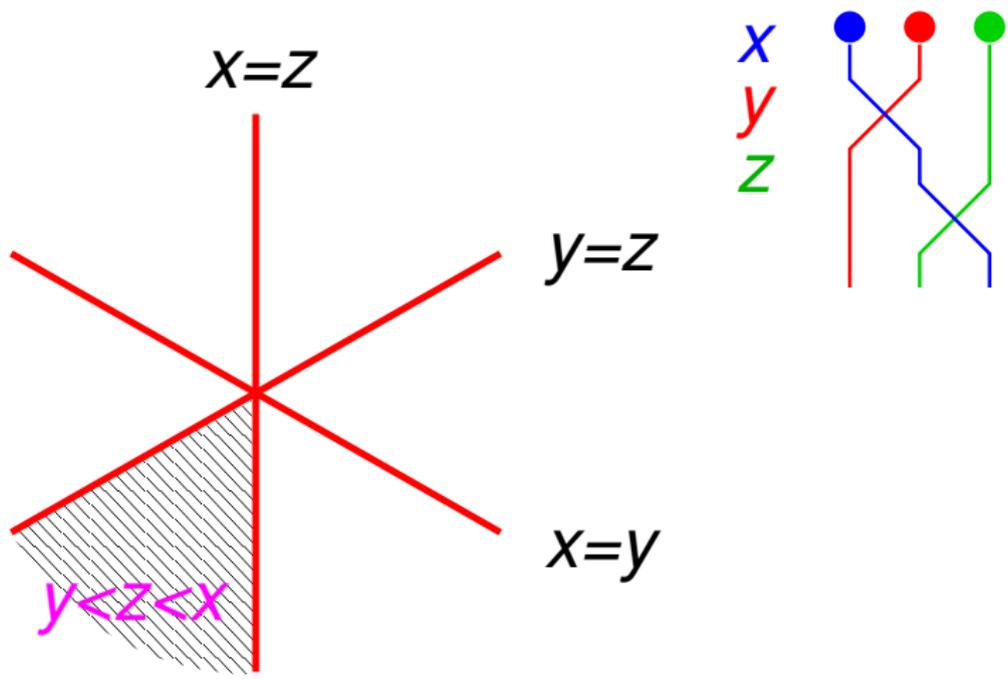
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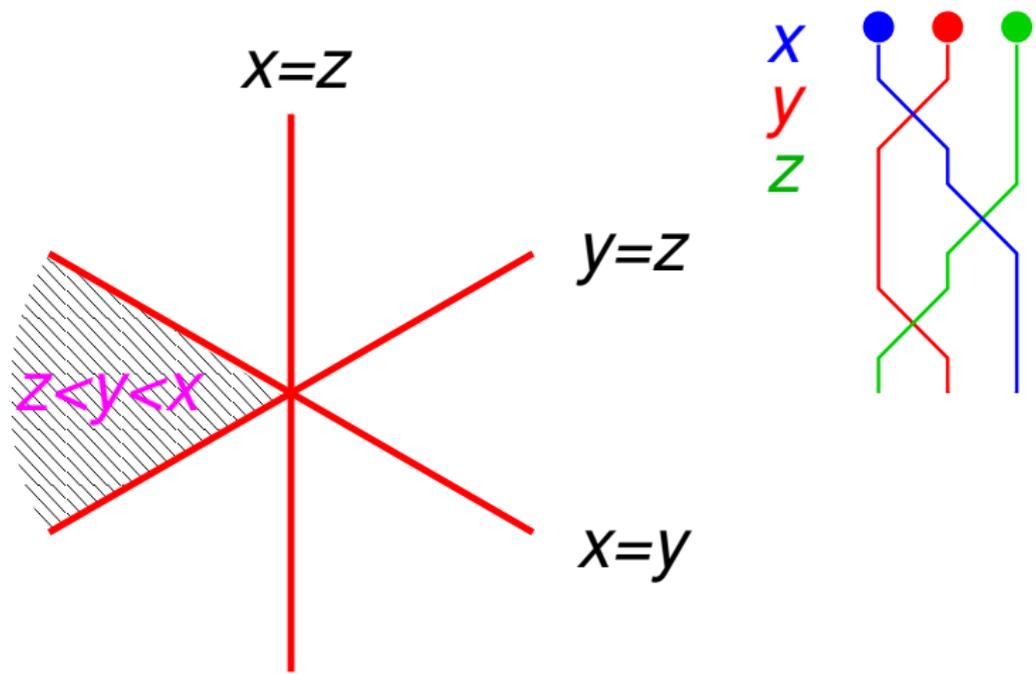
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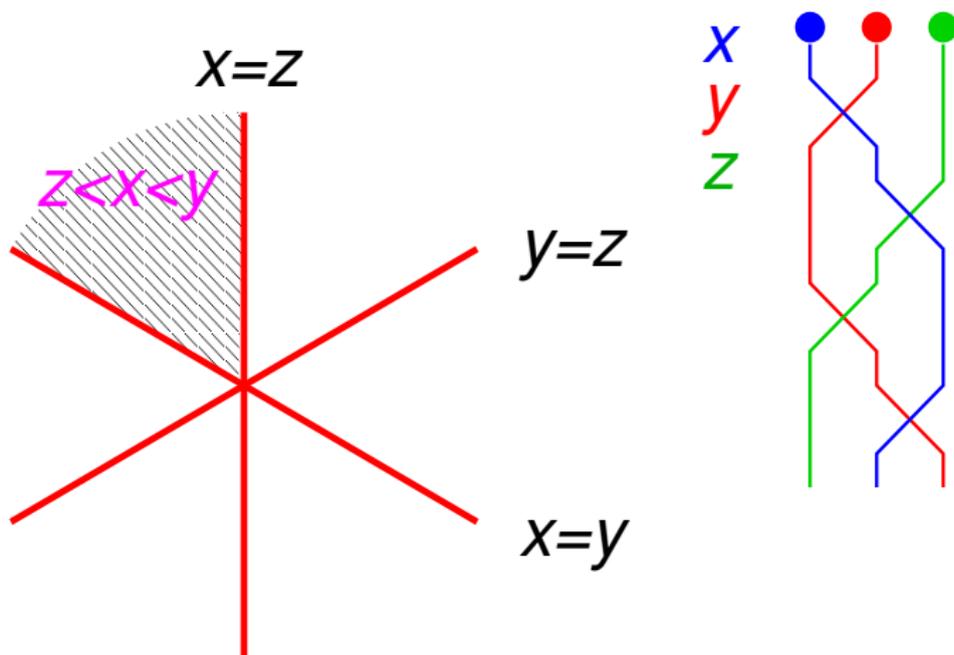
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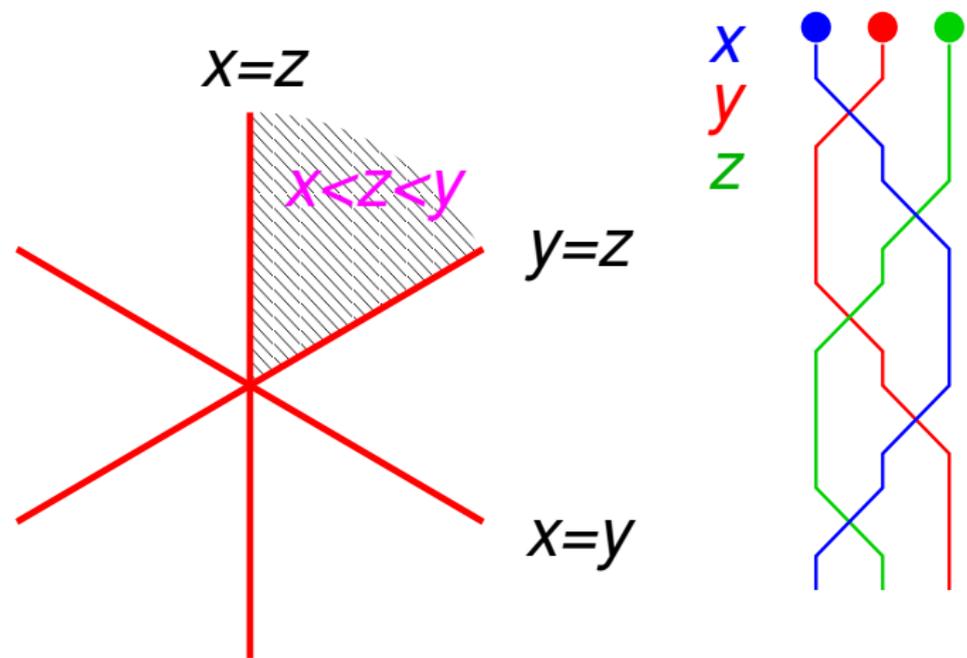
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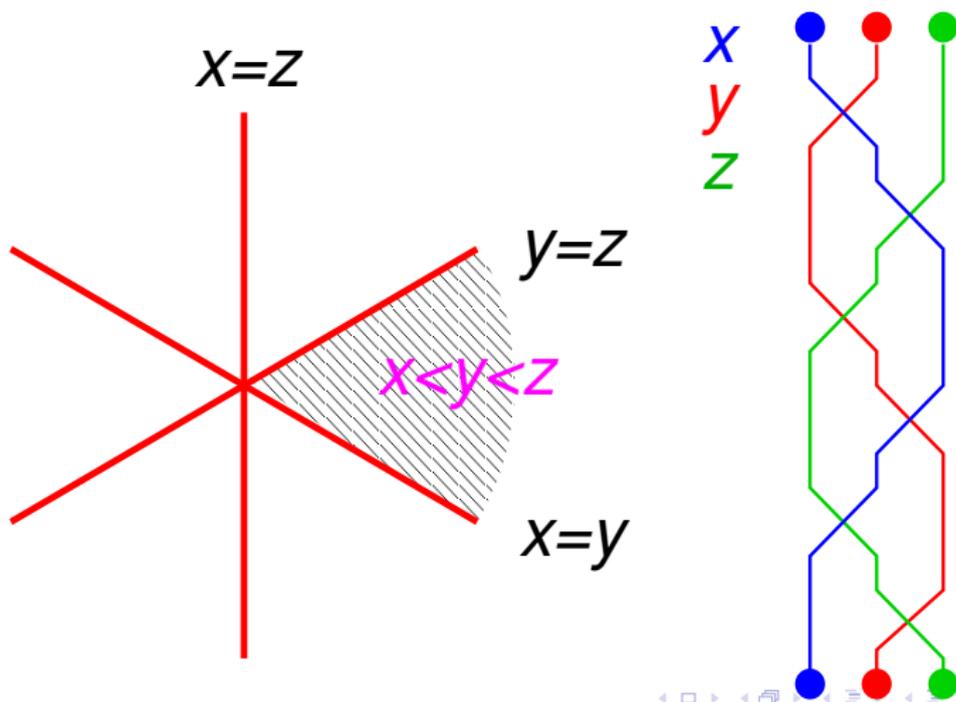
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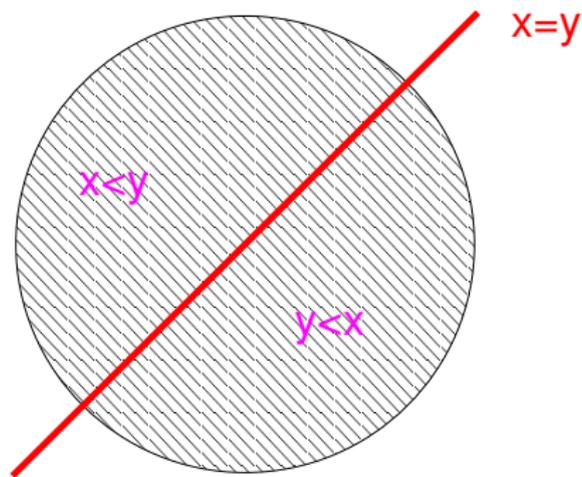


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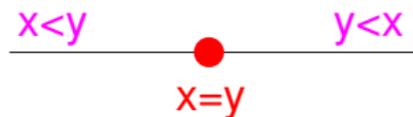
Crossing a border corresponds to reversing one inequality.



# The 2-Dimensional Braid Arrangement



**Braid2**



Projection into 1D

Note that there are 2 regions.

# The 4-Dimensional Braid Arrangement

The arrangement **Braid4** consists of the hyperplanes defined by the equations

$$w = x, \quad w = y, \quad w = z, \quad x = y, \quad x = z, \quad y = z$$

in four-dimensional space.

**Key observation:** We can project **Braid2** from 2D to 1D, and **Braid3** from 3D to 2D,

so, **by analogy**, we should be able to project **Braid4** from 4D to 3D!

## A Technical Interlude

The six equations

$$w = x, \quad w = y, \quad w = z, \quad x = y, \quad x = z, \quad y = z$$

are all satisfied if  $w = x = y = z$ .

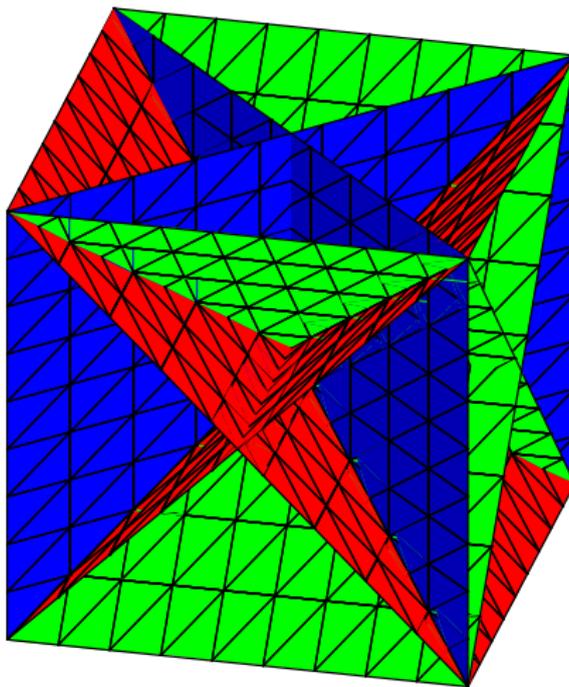
That is, the six hyperplanes of **Braid4** intersect in a common line.

As in the previous cases, we can “squash” (or project) 4D along this line to reduce to 3D.

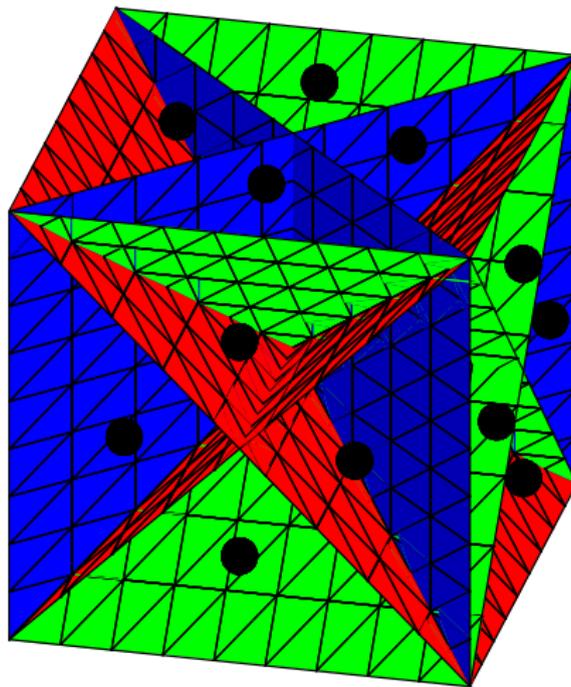
The hyperplane “perpendicular” to that line is defined by  $w + x + y + z = 0$ .

To make the pictures that follow, I gave my computer the equations for **Braid4** and added the equation  $w + x + y + z = 0$ , which means  $w = -x - y - z$ .

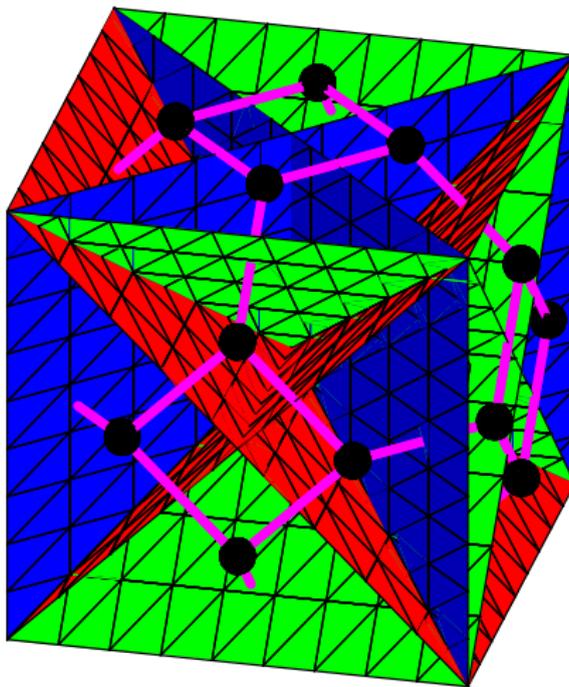
Here's what **Braid4** looks like!



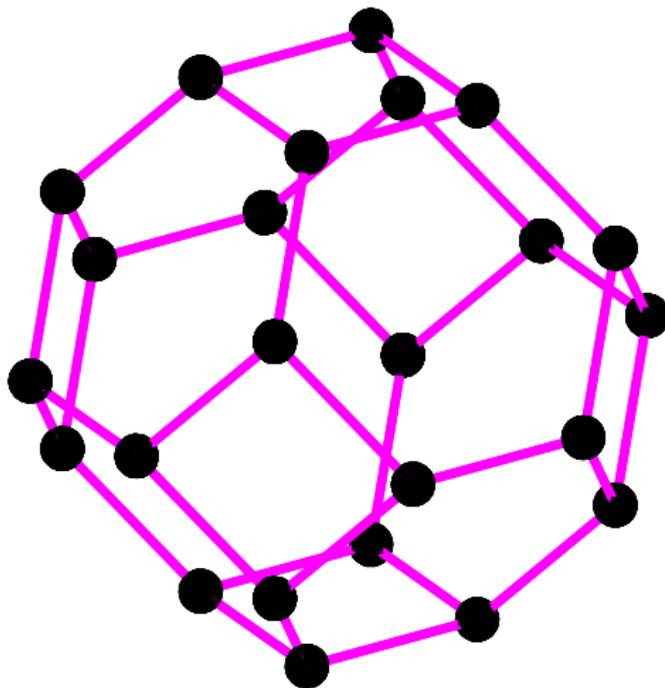
Suppose we put a dot in each region and connect adjacent dots. . .



Suppose we put a dot in each region and connect adjacent dots. . .



... and then remove the hyperplanes, leaving only the dots.



# Regions of Braid<sub>4</sub>

The regions of **Braid<sub>4</sub>** correspond to the orderings of the four coordinates  $w, x, y, z$ :

$wxyz$	$wxzy$	$wyxz$	$wyzx$	$wzxy$	$wzyx$
$xwyz$	$xwzy$	$xywz$	$xyzw$	$xzwy$	$xzyw$
$ywxz$	$ywzx$	$yxwz$	$yxzw$	$yzwx$	$yzxw$
$zwx y$	$zwyx$	$zxwy$	$zxyw$	$zywx$	$zyxw$

- ▶ There are 4 possibilities for the first letter;
- ▶ 3 possibilities for the second, once the first is determined;
- ▶ 2 possibilities for the third, once the first two are determined;
- ▶ only 1 possibility for the last letter.

Total:  $4 \times 3 \times 2 \times 1 = 24$  orderings = 24 regions.

## Regions of Braid4

- ▶ We have just seen that **Braid4** has 24 regions.
- ▶ The regions correspond to permutations of  $w, x, y, z$ .
- ▶ Each region has exactly 3 neighbors.
- ▶ If two regions are adjacent, the corresponding permutations differ by a single flip:

$$\boxed{x \ z \ w \ y} \longleftrightarrow \boxed{x \ w \ z \ y}$$

## Beyond the Fourth Dimension

The  $n$ -dimensional braid arrangement consists of the hyperplanes defined by the equations

$$x_1 = x_2,$$

$$x_1 = x_3, \quad x_2 = x_3,$$

...

$$x_1 = x_n, \quad x_2 = x_n, \quad \dots, \quad x_{n-1} = x_n$$

- ▶ There are  $n(n-1)/2$  hyperplanes (by the staircase formula!)
- ▶ The regions correspond to the possible orderings of the coordinates  $x_1, \dots, x_n$ .
- ▶ The number of regions is  $n \times (n-1) \times \dots \times 3 \times 2 \times 1$  (also known as  $n$  factorial; notation:  $n!$ ).
- ▶ Each region has  $n-1$  neighboring regions.

# Part 3: Cars, Trees, and Scorekeeping

# Parking Cars

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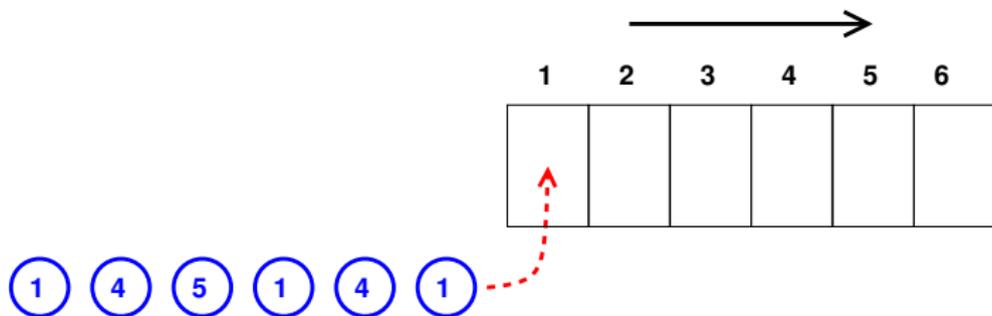
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- ▶ A **parking function** is a list of preferences that allows all cars to park.

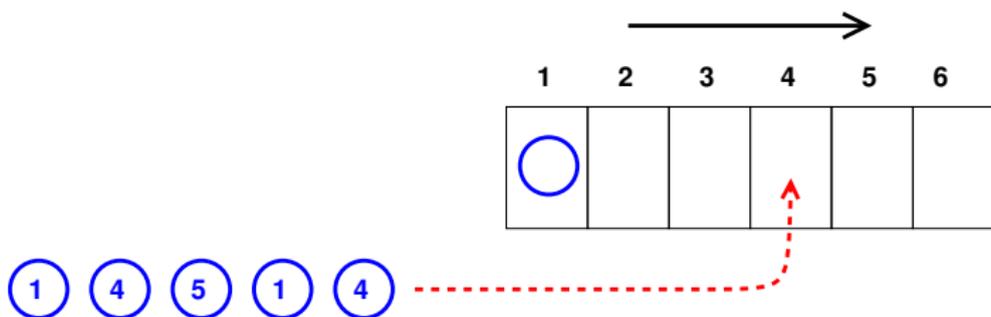
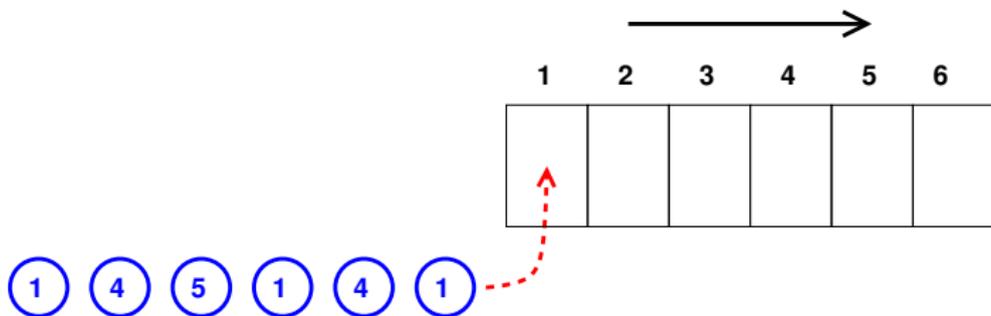
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- ▶ Application: database indexing, hash tables)

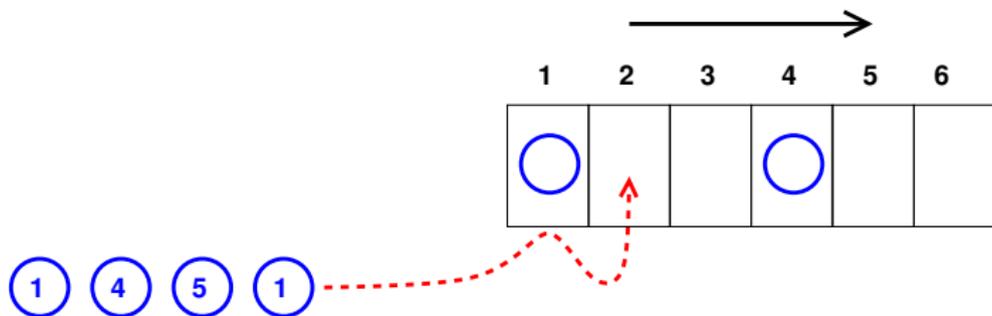
# Parking Functions



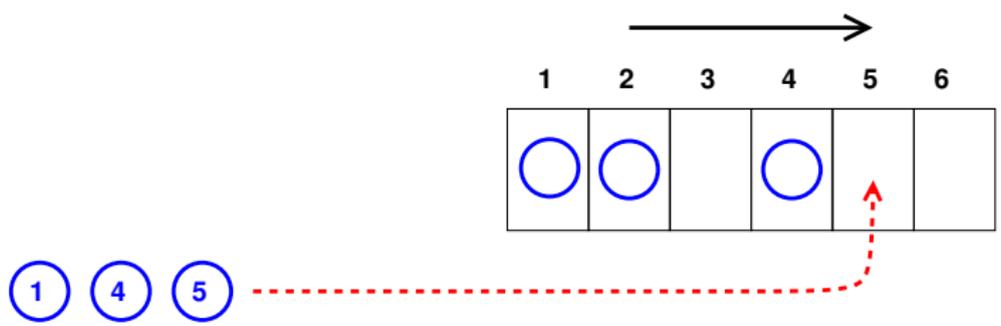
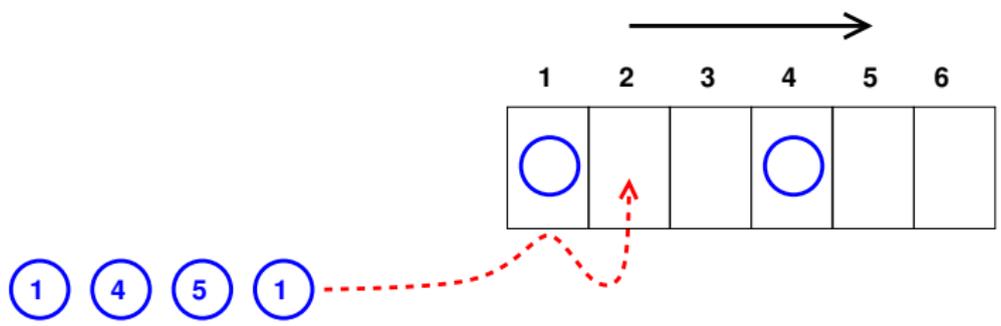
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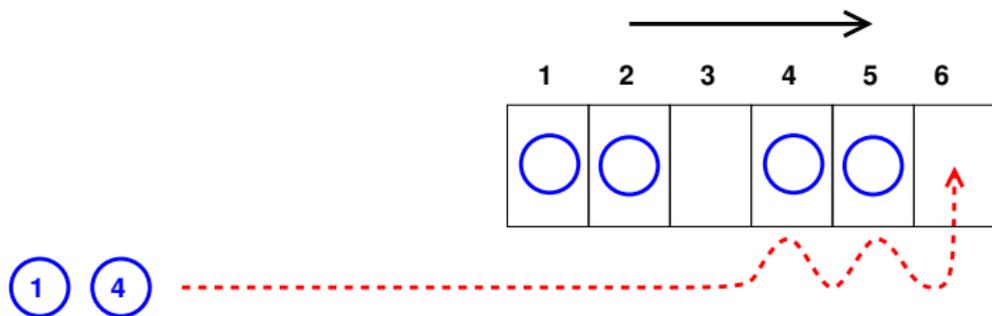
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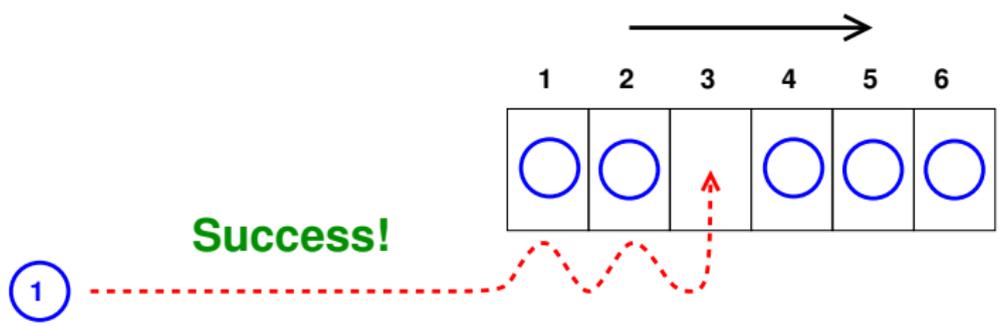
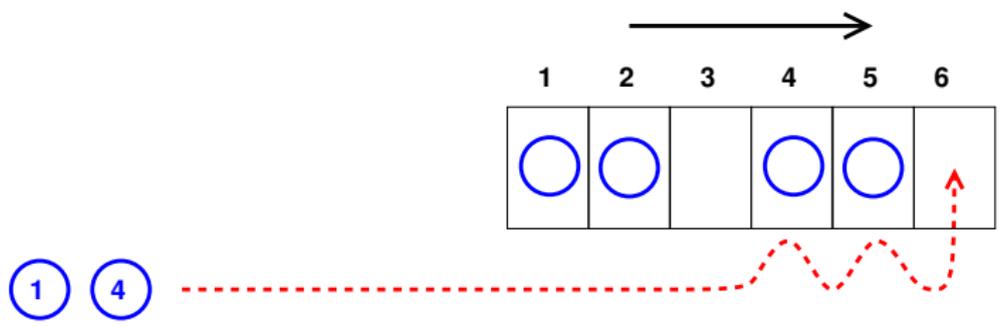
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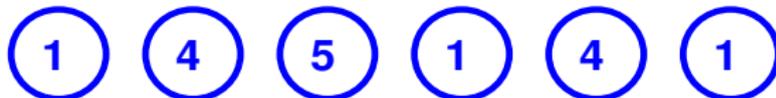


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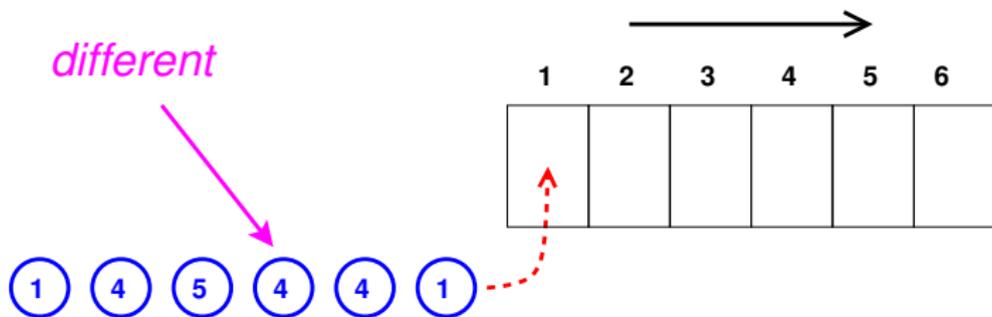
Therefore



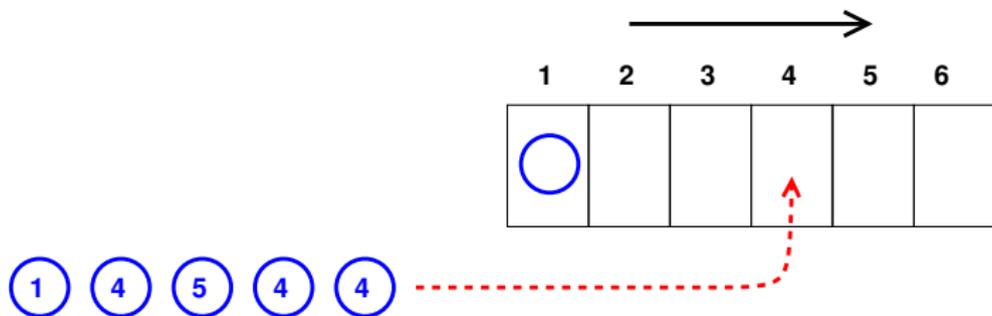
is a parking function. What about



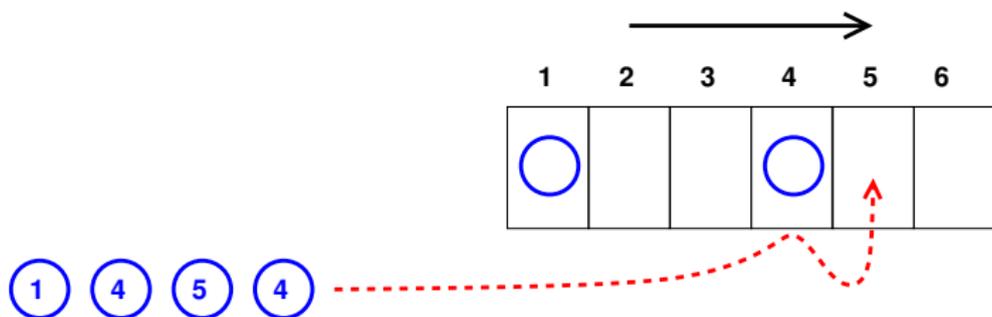
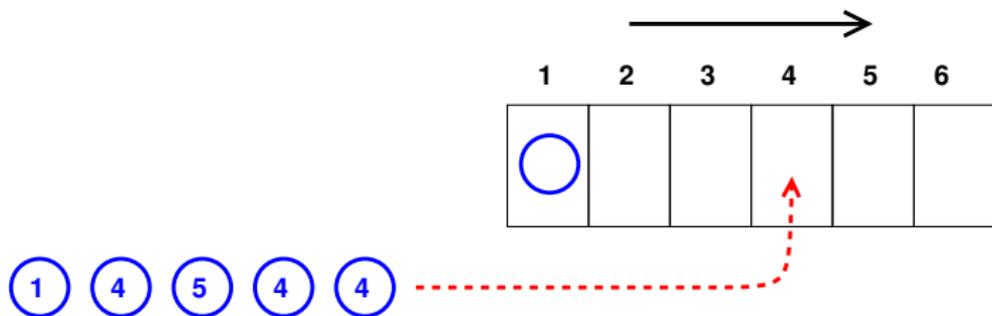
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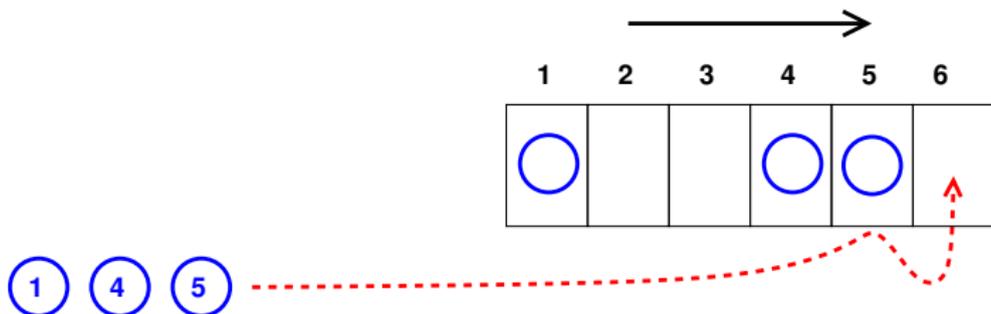
# Parking Functions



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## Parking Two Cars

There are  $4 = 2^2$  possible lists of preferred spots.  
3 of them successfully park both cars.

(1) (1) OK

(2) (1) OK

(1) (2) OK

(2) (2) Not OK

## Parking Three Cars

There are  $27 = 3^3$  possible lists of preferred spots.

16 of them successfully park all three cars.

**Parking functions (the ones that work):**

111	112	122	113	123 132
	121	212	131	213 231
	211	221	311	312 321

**Non-parking functions (the ones that don't work):**

133	222	223	233	333
313		232	323	
331		322	332	

## Parking $n$ Cars

**Observation #1:** Whether or not all the cars can park depends on what their preferred spaces are, but **not on the order** in which they enter the parking lot.

For example, if there are 6 cars and the preference list includes two 5's and one 6, not all cars will be able to park.

Also, every parking function must include at least one 1. (What are some other conditions that must be satisfied?)

## Parking $n$ Cars

**Observation #2:** 3 cars  $\implies$  16 parking functions.

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Number of cars ( $n$ )	Number of parking functions
1	1
2	3
3	16

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**Observation #2:** 3 cars  $\implies$  16 parking functions.

Number of cars ( $n$ )	Number of parking functions
1	1
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3	16
4	125

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**Observation #2:** 3 cars  $\implies$  16 parking functions.

Number of cars ( $n$ )	Number of parking functions
1	1
2	3
3	16
4	125
5	1296

Do you see the pattern?

# Parking $n$ Cars

**Observation #2:** 3 cars  $\implies$  16 parking functions.

Number of cars ( $n$ )	Number of parking functions	
1	1	$= 2^0$
2	3	$= 3^1$
3	16	$= 4^2$
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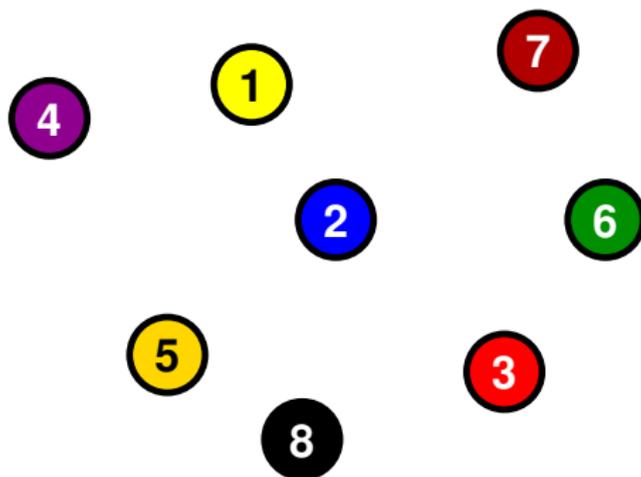
**Conjecture:**  $n$  cars  $\implies (n + 1)^{n-1}$  parking functions.

# Connecting Points

**Problem:** Connect  $n$  points with as few links as possible.

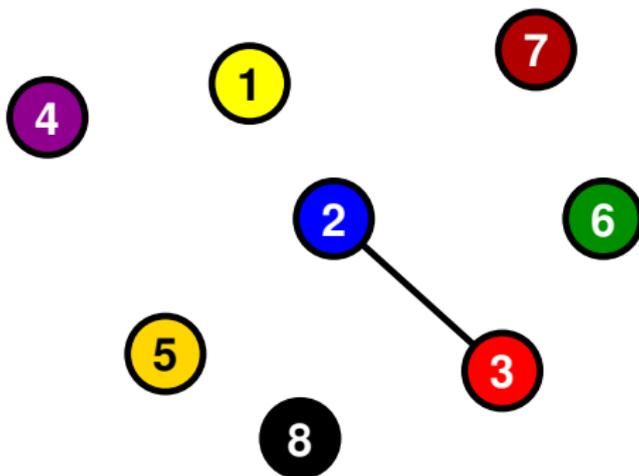
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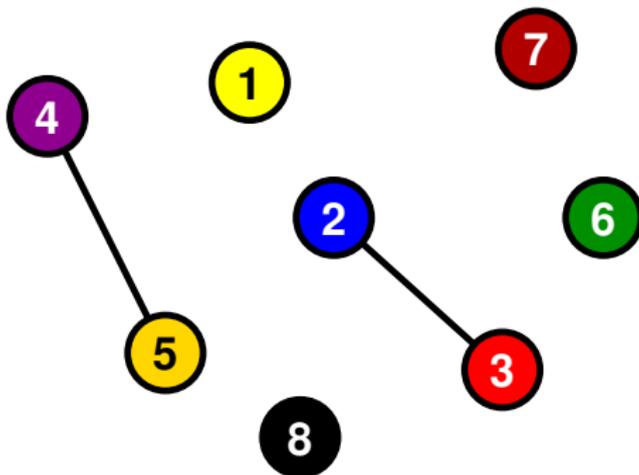
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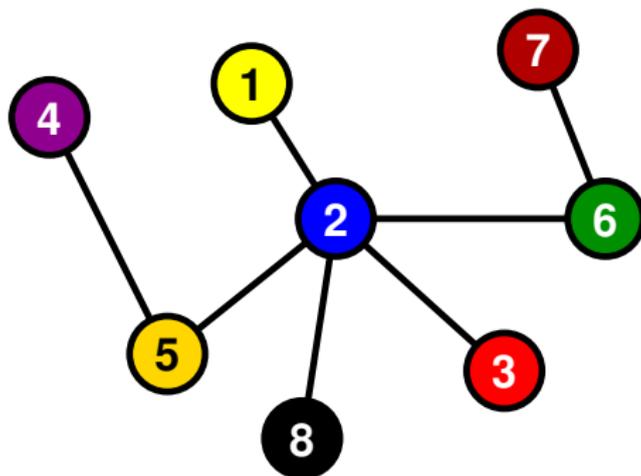
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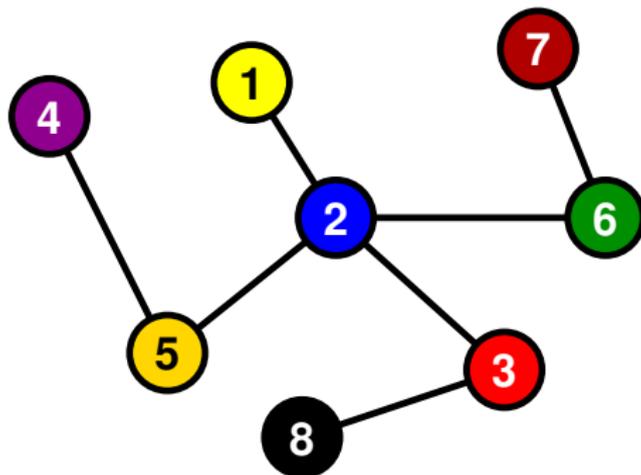
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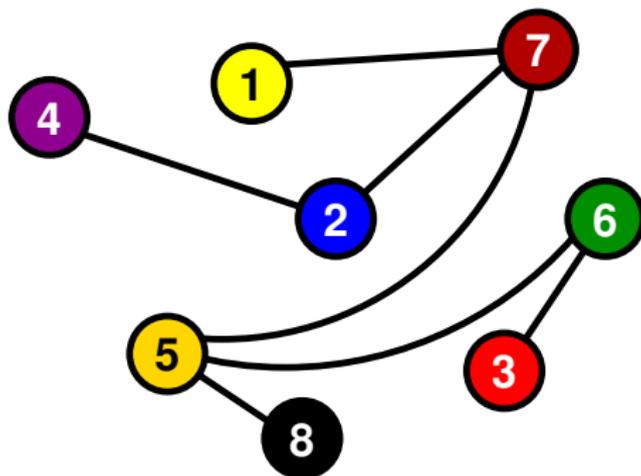
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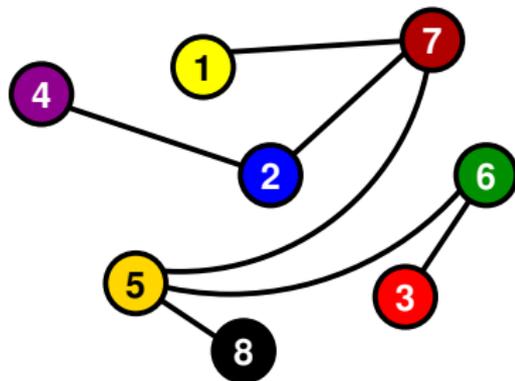
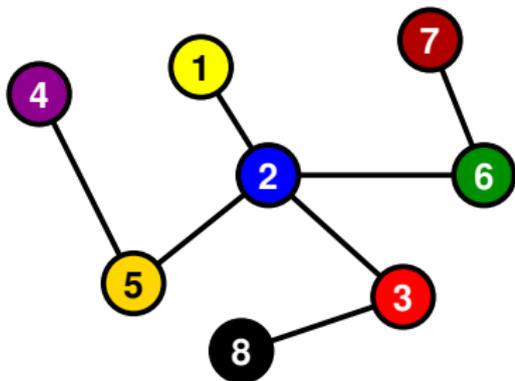
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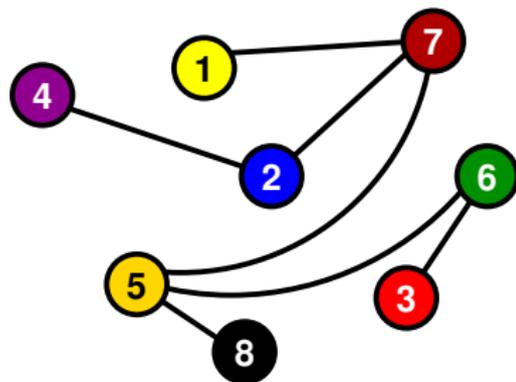
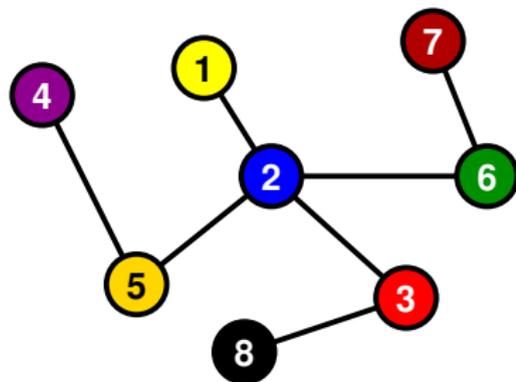
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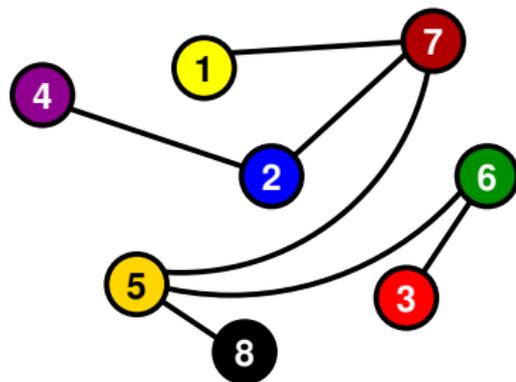
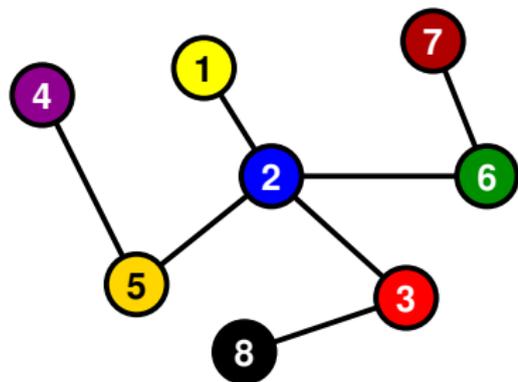
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- ▶ It doesn't matter where the points are or how you draw the links — just which pairs of points are linked.
- ▶ These structures are called **trees**.

# How Many Trees?

1 point:

①

1 tree

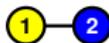
# How Many Trees?

1 point:



1 tree

2 points:



1 tree

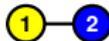
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3 points:



3 trees

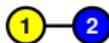
# How Many Trees?

1 point:



1 tree

2 points:



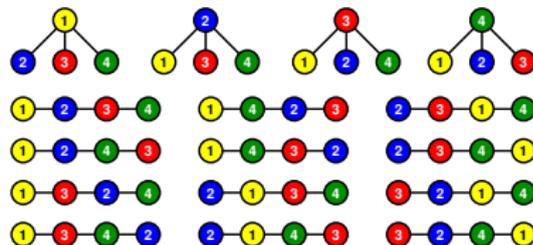
1 tree

3 points:



3 trees

4 points:



16 trees

## Trees and Cars

# Cars	# Parking Functions	# Points	# Trees
1	1	1	1
2	3	2	1
3	16	3	36
4	125	4	16
5	1296	5	125
...		...	
$N$	$(N + 1)^{N-1}$	$N$	$N^{N-2}$

# The Shi Arrangement

The  $n$ -dimensional Shi arrangement consists of the  $n(n-1)$  hyperplanes defined by the equations

$$x_1 = x_2,$$

$$x_1 = x_2 + 1,$$

$$x_1 = x_3,$$

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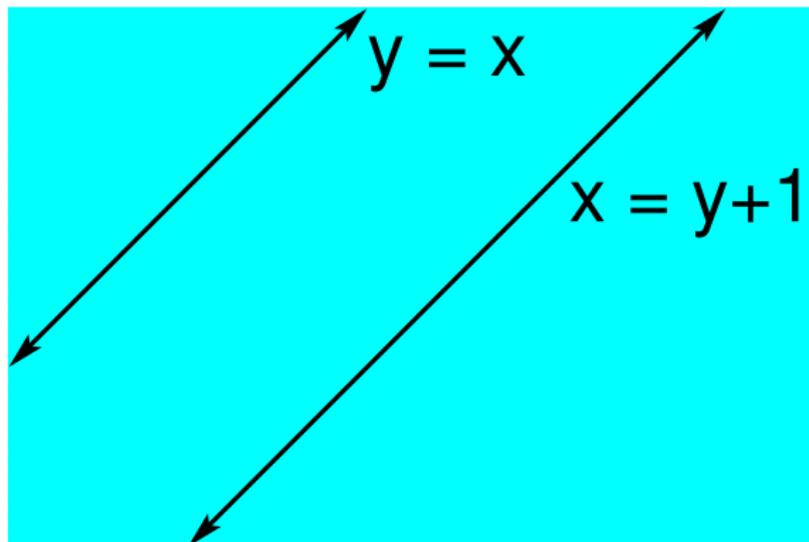
...

$$x_{n-1} = x_n,$$

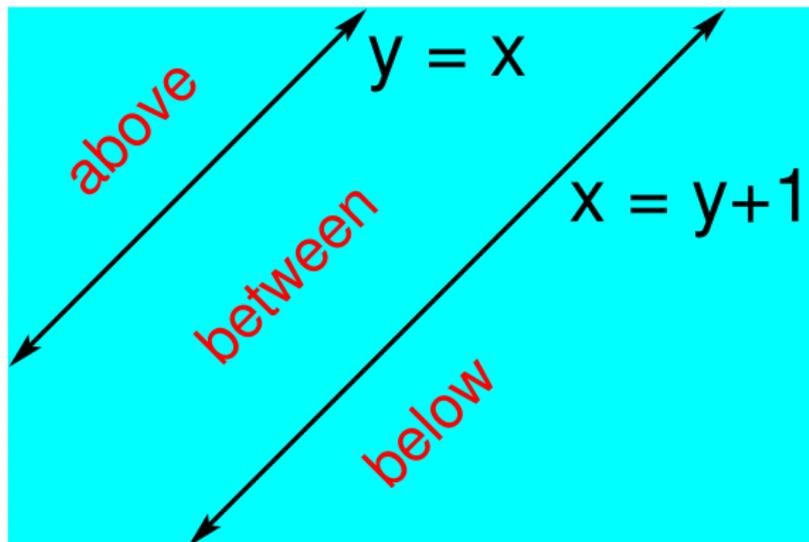
$$x_{n-1} = x_n + 1.$$

(“Take the braid arrangement, make a copy of it, and push the copy a little bit.”)

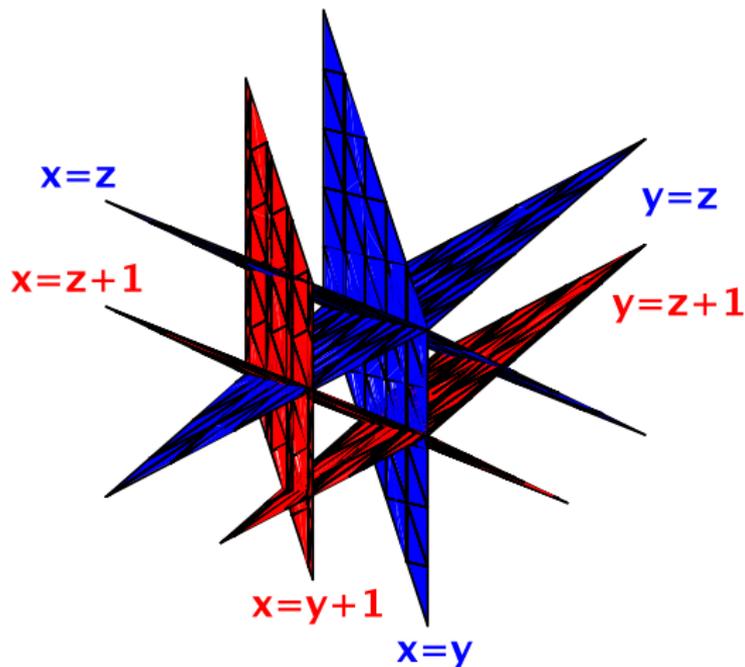
# The 2D Shi Arrangement



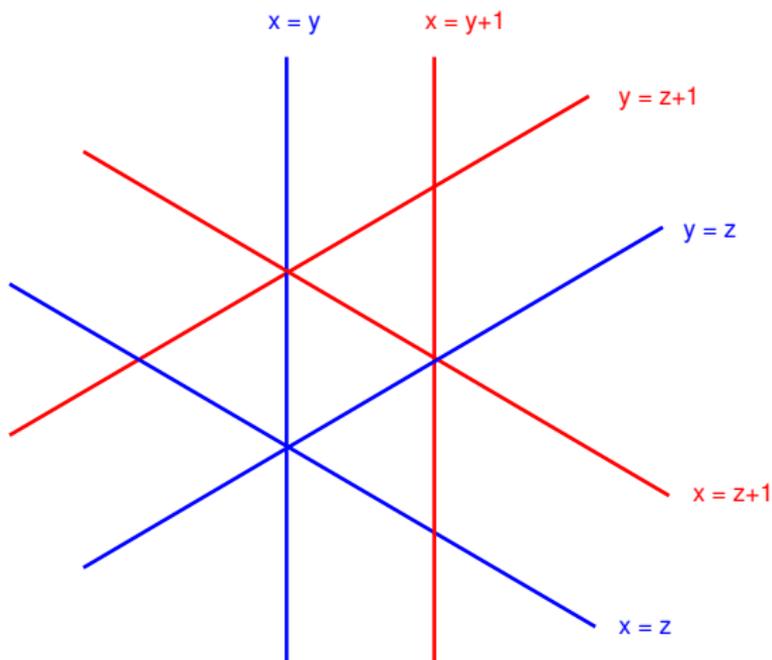
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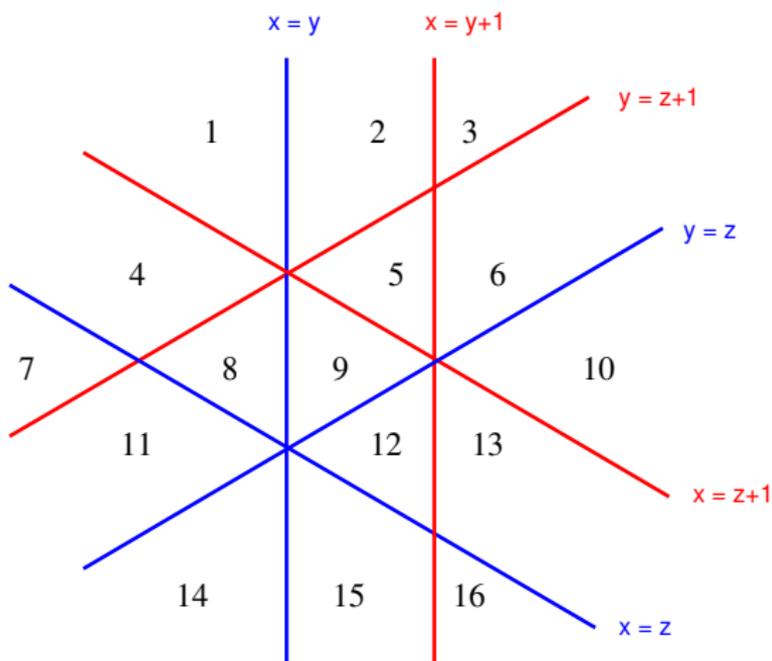
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- ▶ But, in order to score a point against a lower-ranked runner, you must beat him/her by **at least one minute**.
- ▶ **The possible outcomes correspond to regions of the Shi arrangement!**

# Slicing $n$ -Dimensional Space

$(n + 1)^{n-1}$  = number of regions of the Shi arrangement

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$$(n + 1)^{n-1} = \text{number of regions of the Shi arrangement}$$
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## Slicing $n$ -Dimensional Space

$$\begin{aligned}(n + 1)^{n-1} &= \text{number of regions of the Shi arrangement} \\ &= \text{number of handicapped-scoring outcomes} \\ &= \text{number of trees on } n + 1 \text{ points} \\ &= \text{number of ways to park } n \text{ cars}\end{aligned}$$

**Why are all these numbers the same?**

The next figure shows the correspondence between Shi-arrangement regions and parking functions for  $n = 3$ .

