## Scarf complex of powers of an extremal ideal

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#### Example

Let I = (xy, yz, zu), a monomial ideal in S = k[x, y, z, u]. A free resolution of S/I over S:

$$0 \longrightarrow \begin{array}{c} S(xyzu) & \begin{bmatrix} u \\ x \\ -1 \end{bmatrix} & S(xyz) \\ \beta_{3}=1 & \xrightarrow{\oplus} \\ S(xyzu) & \xrightarrow{\oplus} \\ \beta_{2}=3 & & \beta_{1}=3 \end{array} \xrightarrow{\begin{array}{c} z & 0 & zu \\ -x & u & 0 \\ 0 & -y & -xy \end{bmatrix}} & \begin{array}{c} S(xy) \\ \oplus \\ S(yz) \\ \oplus \\ S(zu) \\ \beta_{1}=3 \end{array}$$

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# Taylor complex [Taylor '66]





faces = lcm

always a simplex

The Taylor complex supports the free resolution of any monomial ideal (but not minimally)

$$0 \longrightarrow S(xyzu) \xrightarrow{\begin{pmatrix} u \\ x \\ -1 \end{pmatrix}} \begin{array}{c} S(xyz) \\ \oplus \\ S(yzu) \\ \oplus \\ S(xyzu) \end{array} \xrightarrow{\begin{pmatrix} z & 0 & zu \\ -x & u & 0 \\ 0 & -y & -xy \end{pmatrix}} \begin{array}{c} S(xy) \\ \oplus \\ S(yz) \\ \oplus \\ S(zu) \end{array} \xrightarrow{\oplus} \begin{array}{c} S(zu) \\ \oplus \\ S(zu) \end{array}$$

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# Scarf complex [BPS '98]



Scarf complex removes all matching faces with same label.

- It might not support a free resolution ....
- but it is a lower bound:
- Up to isomorphisms, any free resolution of monomial ideal *I* contains Scarf(*I*) as a subcomplex.

$$0 \longrightarrow S(xyz) \oplus S(yzu) \xrightarrow{\begin{bmatrix} z & 0 \\ -x & u \\ 0 & -y \end{bmatrix}} S(xy) \oplus S(yz) \oplus S(zu) \longrightarrow S$$
$$\beta_2 = 2 \qquad \qquad \beta_1 = 3$$

Power of a (monomial) ideal:

$$I^2 = (m_1, \ldots, m_p)^2 = (\{m_i m_j \colon 1 \le i \le j \le p\}).$$

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Similarly for  $I^r$  when r > 2. Taylor is never minimal for  $I^r$  when  $r \ge 2$ .

# Extremal ideal [CEFMM\$S '24]

$$S_{\mathcal{E}} = k[x_A \colon \emptyset \neq A \subseteq [q]]$$
  
 $\epsilon_i = \prod_{i \in A \subseteq [q]} x_A$ 

Then  $\mathcal{E}_q = (\epsilon_1, \ldots, \epsilon_q)$  is the extremal ideal

$$\mathcal{E}_3 = (\epsilon_1 = x_1 x_{12} x_{13} x_{123}, \epsilon_2 = x_2 x_{12} x_{23} x_{123}, \epsilon_3 = x_3 x_{13} x_{23} x_{123})$$

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$$\beta_i(I^r) \leq \beta_i(\mathcal{E}_q^r)$$

where *I* is any ideal generated by *q* square-free monomials. So we focus on  $\mathbb{S}_q^r = \text{Scarf}(\mathcal{E}_q^r)$ .

- What does it look like?
- Does it always support a minimal free resolution of \mathcal{E}\_a^r?

Theorem (EF\$S '24) Let  $\sigma = \{ \boldsymbol{\epsilon}^{\mathbf{a}_1}, \dots, \boldsymbol{\epsilon}^{\mathbf{a}_d} \} \in \mathbb{T}_a^r = Taylor(\mathcal{E}_a^r)$ . Then  $\sigma \in \mathbb{S}_a^r$  iff •  $\sigma' \in \mathbb{S}^{r}_{\sigma}$  for all proper subsets  $\sigma'$  of  $\sigma$ ; and **b** =  $\mathbf{a}_i$  are only solutions  $\mathbf{b} \in \mathcal{N}_q^r = {\mathbf{c} \in \mathbb{N}^q : |\mathbf{c}| = r}$  to:  $\mathbf{b} \cdot \mathbf{e}_A \le \max{\{\mathbf{a}_i \cdot \mathbf{e}_A\}}$  for all  $A \subseteq [q]$ where  $\mathbf{e}_A = \sum_{i \in A} \mathbf{e}_i$ (0, 0, 2)(1, 0, 1)(0, 1, 1)(0, 2, 0)(2, 0, 0)(1, 1, 0)< ロ > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ >

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## What does $\mathbb{S}_{a}^{r}$ look like: geometric simplifications

Whether or not a set of vertices is a face of  $\mathbb{S}_q^r$  is invariant under:

- translation (subtract a common vector)
- permutation of coordinates
- truncation of common 0's (more generally, common entries)

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 $\{3000, 1011\}$ 

- translation (subtract a common vector)
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 $\{3000,1011\} \Leftrightarrow \{300,111\} \Leftrightarrow \{200,011\} \Leftrightarrow \{002,110\} \text{ no}$ 

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Checking edges reduces to checking  $\{0u, v0\}$  where u, v are partitions. This allows us to find, by computer search, all edges for  $r \leq 8$  and arbitrary q.

 $q \leq 4$ 



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$$U_q^r = \{v \in \{0,1\}^q : |v| = r\}$$

- Always a facet
- All facets are translations of  $U_q^r$
- Only facet containing  $(r, 0, \ldots, 0)$  is  $(r 1, 0, \ldots, 0) + U_q^1$
- For q = 4, it is drawn in 3-dimensions (octahedron), but it is actually 5-dimensional simplex with 6 vertices. (And similarly for larger q.)



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#### $\{21000, 00111, 11100, 11010, 11001, 10110, 10101, 10011\}$

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Minimal non-faces (up to permutation, and padding with 0's)

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Minimal non-faces (up to permutation, and padding with 0's) {30,03}, {30,12}, {300,021}, {300,111}, {3000,0111}

{21000,00111,11100,11010,11001,10110,10101,10011}

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Minimal non-faces (up to permutation, and padding with 0's)

 $\blacktriangleright \ \{30,03\}, \{30,12\}, \{300,021\}, \{300,111\}, \{3000,0111\}$ 

 $\blacktriangleright \ \{210,021\},\{210,012\},\{2100,0111\}$ 

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Reflection is also a rigid motion, and so preserves being a face or not.

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## Example (Reflecting through the origin)

 $\{3000, 2100, 2010, 2001\}$  is a face, so its negative (reflecting through the origin) satisfies all the conditions except for being non-negative.

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## Example (Reflecting through the origin)

{3000, 2100, 2010, 2001} is a face, so its negative (reflecting through the origin) satisfies all the conditions except for being non-negative. Fix this by then also translating by adding the max of each original coordinate, 3111. This gives {0111, 1011, 1101, 1110} which is a face.

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{3000, 2100, 2010, 2001} is a face, so its negative (reflecting through the origin) satisfies all the conditions except for being non-negative. Fix this by then also translating by adding the max of each original coordinate, 3111. This gives {0111, 1011, 1101, 1110} which is a face.

We can also reflect through a plane, again translating afterwards to restore non-negativity. But in some cases, it may be hard to stay integral.

## Theorem (EF\$S '24)

 $\mathbb{S}_q^r$  supports a minimal free resolution of  $\mathcal{E}_q^r$  for r = 2 and for  $q \leq 4$ .

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## Theorem (new)

 $\mathbb{S}_q^3$  supports a minimal free resolution of  $\mathcal{E}_q^3$ . In particular, it supports a free resolution of  $I^3$  for any monomial ideal with q generators.

## Proof.

We find a homogeneous acyclic matching [BW '02] (from discrete Morse theory) that leaves exactly the Scarf complex.

# Homogeneous acyclic matching [BW '02]



- Directed graph of poset of faces of Taylor complex, labeled by lcm
- Partial matching *M* using only poset-edges whose poset-vertices have the same label
- $\blacktriangleright$  Reverse arrows on poset-edges of  ${\cal M}$
- If new poset is acyclic, then removing *M* leaves a complex supporting free resolution

**Scarf:** 
$$f_{i-1}(\mathbb{S}_q^3) = \binom{\binom{q}{3}}{i} + \text{lower order terms}$$

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**Scarf:** 
$$f_{i-1}(\mathbb{S}_q^3) = \begin{pmatrix} \binom{q}{3} \\ i \end{pmatrix} + \text{lower order terms}$$

Dominant term comes from U<sup>3</sup><sub>q</sub>

- Question: For arbitrary r, is  $\binom{\binom{q}{r}}{r}$  always dominant?
- This provides an upper bound for the size of a free resolution of I<sup>3</sup> for any ideal I.

Taylor: 
$$f_{i-1}(\mathbb{T}(\mathcal{E}_q^3)) = \begin{pmatrix} \binom{q+2}{3} \\ i \end{pmatrix}$$

Both bounds are  $O(q^{3i})$ , but ...

## Comparing upper bounds



# Comparing upper bounds



For large q

$$\max_{i} f_{i}(\mathbb{T}_{q}^{3}) / \max_{i} f_{i}(\mathbb{S}_{q}^{3}) = 2^{q^{2} + O(q)} / \sqrt{1 + 6/q}$$

## Comparing upper bounds



For large q; and fixed  $0 \le c < 1$ 

$$\max_{i} f_{i}(\mathbb{T}_{q}^{3}) / \max_{i} f_{i}(\mathbb{S}_{q}^{3}) = 2^{q^{2} + O(q)} / \sqrt{1 + 6/q}$$
$$f_{i-1}(\mathbb{T}_{q}^{3}) / f_{i-1}(\mathbb{S}_{q}^{3}) = (1 - c)^{-q^{2} + O(q)} O(e^{-3cq/(1 - c)})$$

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where  $i = c(\dim(\mathbb{S}_q^3)) + 1$ .

Describe Scarf complex (facets and/or minimal non-faces)

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Prove Scarf complex supports resolution

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Prove Scarf complex supports resolution

#### general r (for you)

- Describe Scarf complex (facets and/or minimal non-faces)
- Prove Scarf complex supports resolution
- general r (for you)
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Prove Scarf complex supports resolution

- r = 4 (for you)
  - Describe Scarf complex (facets and/or minimal non-faces)
  - Prove Scarf complex supports resolution
- general r (for you)
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  - Prove Scarf complex supports resolution

Further properties of extremal ideal (we are working on this)

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